# Introduction to Optimization

### **Outline:**

- Standard form optimization problem and terminology.
- Convex optimization problems.
- Lagrange duality.
- Optimization methods.

Optimization problem in standard form

minimize 
$$f_0(\boldsymbol{x})$$
 subject to  $f_i(\boldsymbol{x}) \leq 0$  ,  $i=1,\ldots,m$   $h_i(\boldsymbol{x}) = 0$  ,  $i=1,\ldots,p$ 

where

 $x \in \mathbb{R}^n$  is the **optimized vector** of variables.

 $f_0: \mathbb{R}^n \to \mathbb{R}$  is the **objective function**.

 $f_i: \mathbb{R}^n \to \mathbb{R}$  is the **inequality constraint** function.

 $h_i: \mathbb{R}^n \to \mathbb{R}$  is the **equality constraint** function.

**Explicit constraints** are  $f_i(x) \le 0$  and  $h_i(x) = 0$ ; unconstrained problem has no explicit constrains (i.e. m = p = 0).

**Implicit constraint** is  $x \in \mathcal{D}$  where  $\mathcal{D}$  is a common domain of the objective function and constraint functions

$$\mathcal{D} = \bigcap_{i=0}^m \mathbf{dom} \, f_i \cap \bigcap_{i=1}^p \mathbf{dom} \, h_i \, .$$

Feasible set: contains points which satisfy implicit and explicit constraints

$$\mathcal{X}_{\text{feas}} = \mathcal{D} \cap \{ \boldsymbol{x} \mid f_i(\boldsymbol{x}) \leq 0, i = 1, ..., m, h_j(\boldsymbol{x}) = 0, j = 1, ..., n \}$$

Example: (minimal entropy discrete distribution)

minimize 
$$-\sum_{i=1}^{n} x_i \log x_i$$
 subject to 
$$\sum_{i=1}^{n} x_i = 1 .$$

which has explicit constraint  $\sum_{i=1}^{n} x_i = 1$ , implicit constraints  $x_i > 0$  and feasible set  $\mathcal{X}_{\text{feas}} = \{ \boldsymbol{x} \mid \sum_{i=1}^{n} x_i = 1, x_i > 0, i = 1, \ldots, n \}$ .

LP problem	QP problem
minimize $oldsymbol{c}^T oldsymbol{x}$ subject to $oldsymbol{\mathbf{A}} oldsymbol{x} = oldsymbol{b}$ $oldsymbol{\mathbf{D}} oldsymbol{x} \leq oldsymbol{q}$	minimize $rac{1}{2}m{x}^T\mathbf{H}m{x}+m{c}^Tm{x}$ subject to $\mathbf{A}m{x}=m{b}$ $\mathbf{D}m{x} \leq m{q}$

where

 $x \in \mathbb{R}^n$  is a vector of optimized variables

 $oldsymbol{c} \in \mathbb{R}^n$ ,  $oldsymbol{b} \in \mathbb{R}^p$ ,  $oldsymbol{q} \in \mathbb{R}^m$  are vectors

 $\mathbf{A} \in \mathbb{R}^{p \times n}$ ,  $\mathbf{D} \in \mathbb{R}^{m \times n}$ ,  $\mathbf{H} \in \mathbb{R}^{n \times n}$  are matrices

Note that LP and QP can be always rewritten to a simpler form using the **slack variables trick:** the inequality constraints  $\mathbf{D}x \leq q$  are replaced by equivalent constraints  $\mathbf{D}x + \xi = q$  and  $\xi \geq 0$ .

# (Globally) optimal value:

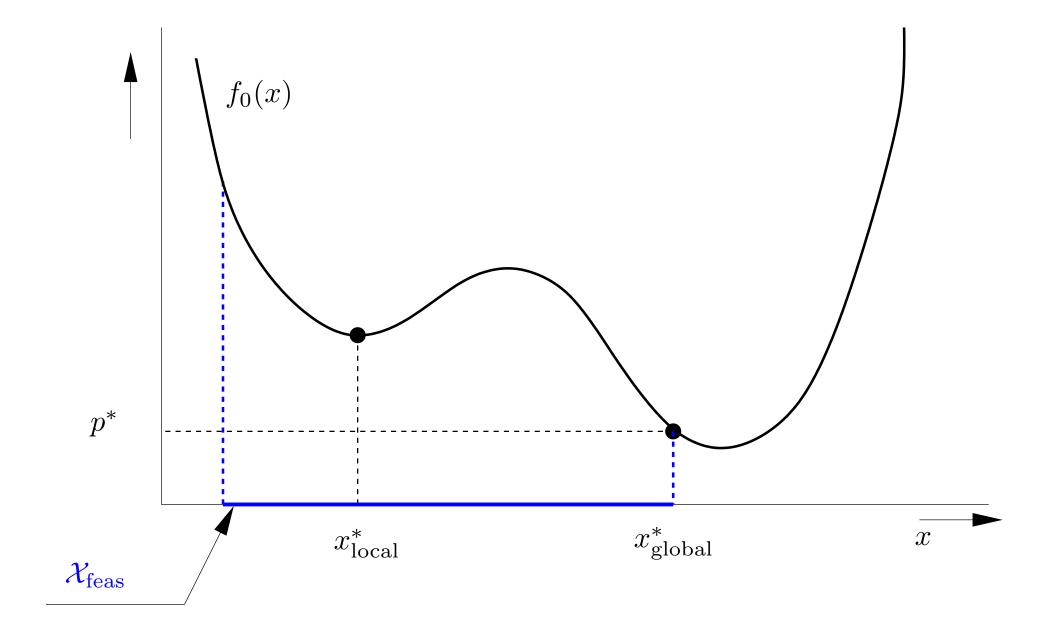
$$p^* = \inf\{f_0(\boldsymbol{x}) \mid \boldsymbol{x} \in \mathcal{X}_{\text{feas}}\}$$

- $p^* = \infty$  if the problem is infeasible, i.e.,  $\mathcal{X}_{\text{feas}} = \{\emptyset\}$ .
- $p^* = -\infty$  if the problem is unbounded.

**Optimal solutions:** x is the optimal solution if it is feasible and  $f(x) = p^*$ ;  $\mathcal{X}_{opt} = \{x \mid f_0(x) = p^*, x \in \mathcal{X}_{feas}\}$  is the set of optimal solutions.

**Locally optimal:**  ${\boldsymbol x}$  is locally optimal if there exist R>0 such that  ${\boldsymbol x}$  is optimal for

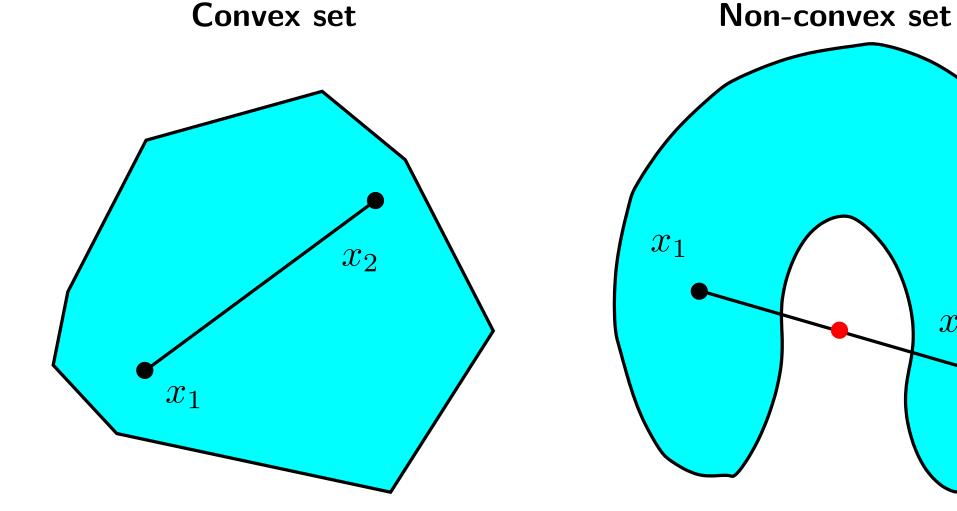
minimize 
$$f_0(\boldsymbol{y})$$
 subject to  $\boldsymbol{y} \in \mathcal{X}_{\mathrm{feas}} \cap \{\boldsymbol{y} \mid \|\boldsymbol{x} - \boldsymbol{y}\| \leq R\}$ 



 $x_2$ 

A set  $\mathcal{X} \subseteq \mathbb{R}^n$  is convex if the line segment connecting any two points from  $\mathcal X$  lies in  $\mathcal X$ , i.e., for all  $x_1,x_2\in\mathcal X$  and all  $\theta$ such that  $0 \le \theta \le 1$  it holds

$$\boldsymbol{x}_1(1-\theta) + \theta \boldsymbol{x}_2 \in \mathcal{X}$$
.



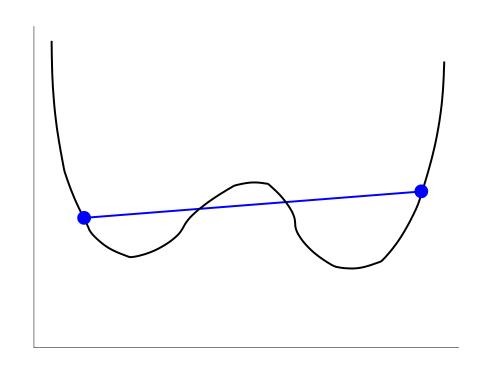
A function  $f \in \mathbb{R}^n \to \mathbb{R}$  is convex if  $\operatorname{dom} f$  is convex and for all  $x_1$ ,  $x_2 \in \operatorname{dom} f$  and all  $\theta$  such that  $0 \le \theta \le 1$  it holds

$$f(x_1(1-\theta) + x_2\theta) \le f(x_1)(1-\theta) + f(x_2)\theta$$
.

### **Convex function**

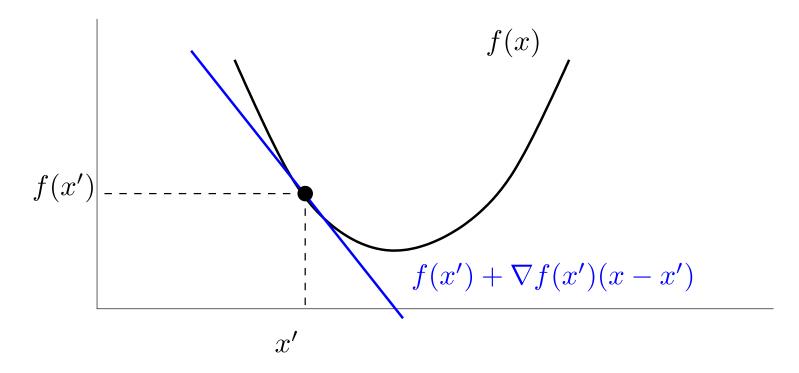
# $f(x_2)$ $f(x_1)(1-\theta) + f(x_2)\theta$ $x_1$

### Non-convex function



**First-order condition:** Suppose that  $f: \mathbb{R}^n \to \mathbb{R}$  is differentiable, i.e., gradient  $\nabla f(x) \in \mathbb{R}^n$  exists at each point  $x \in \operatorname{dom} f$ . Then f is convex if and only if  $\operatorname{dom} f$  is convex and

$$f(\boldsymbol{x}) \geq f(\boldsymbol{x}') + \nabla f(\boldsymbol{x}')^T(\boldsymbol{x} - \boldsymbol{x}') , \qquad \forall \boldsymbol{x}, \boldsymbol{x}' \in \mathcal{X} .$$



**Second-order condition:** Suppose that f twice differentiable, i.e., the Hessian matrix of second derivatives  $\nabla^2 f(x)$  exists at each point  $x \in \operatorname{dom} f$ . Then f is convex if and only if  $\operatorname{dom} f$  is convex and  $\nabla^2 f(x)$  is positive semi-definite for all  $x \in \operatorname{dom} f$ .

The optimization problem is convex if the objective function  $f_0(x)$  is convex and the feasible set  $\mathcal{X}_{\text{feas}}$  is convex.

- In particular, the problem is convex if  $f_0, f_1, \ldots, f_m$  are convex and the equality constraints  $h_i$  are affine, i.e.,  $h_i(\mathbf{x}) = \mathbf{a}_i^T \mathbf{x} b_i = 0$ .
- The standard form of the convex optimization problem

```
minimize f_0(m{x}) subject to f_i(m{x}) \leq 0 , i=1,\ldots,m m{A}m{x} = m{b}
```

- Basic property of the convex problems: Any locally optimal solution is globally optimal  $\Rightarrow$  greatly simplifies optimization.
  - We can use **descent methods**: iteratively move in a descent direction until we reach the optimum.
  - For non-convex problems we can get stuck in a local optimum; it is difficult to identify whether the attained optimum is local or global.

# LP problemQP problemminimize $c^Tx$ minimize $\frac{1}{2}x^T\mathbf{H}x + c^Tx$ subject to Ax = bsubject to Ax = b $Dx \le q$ $Dx \le q$

**Linear Programming** is a convex problem since the objective is a convex function, the equality functions are affine, the inequality constraints define a convex set.

**Quadratic Programming** is a convex problem if and only if the matrix  ${\bf H}$  is positively semi-definite;

Recall the Second-order condition and notice that for QP the Hessian matrix  $\nabla^2 f(x) = \mathbf{H}$ .

Suppose that  $f_0$  is differentiable. Then a vector  $m{x}$  is the optimal solution if and only if it is feasible  $m{x} \in \mathcal{X}_{\mathrm{feas}}$  and

$$\nabla f_0(\boldsymbol{x})^T(\boldsymbol{y}-\boldsymbol{x}) \geq 0$$
 for all  $\boldsymbol{y} \in \mathcal{X}_{\mathrm{feas}}$ .

### How to see this?

Recall the definition of the directional derivative

$$f_0'(\boldsymbol{x}; \boldsymbol{\delta}) = \lim_{h \to 0_+} \frac{f_0(\boldsymbol{x} + h\boldsymbol{\delta})}{h} = \nabla f_0(\boldsymbol{x})^T \boldsymbol{\delta}.$$

The sign of  $f_0'(x; \delta)$  determines whether  $f_0$  increases or decreases when we move from x in the direction  $\delta$ .

- Moving from a feasible point x along a feasible direction  $\delta = y x$ ,  $y \in \mathcal{X}_{\mathrm{feas}}$  by sufficiently small step produces a feasible point.
- A vector x is optimal iff there is no feasible direction which decreases the objective function, i.e., for each  $y \in \mathcal{X}_{\mathrm{feas}}$  moving along  $\delta = y x$  increases the objective so that

$$f_0'(\boldsymbol{x};\boldsymbol{\delta}) \geq 0 \quad \Rightarrow \quad \nabla f_0(\boldsymbol{x})^T \boldsymbol{\delta} \geq 0 \quad \Rightarrow \nabla f_0(\boldsymbol{x})^T (\boldsymbol{y} - \boldsymbol{x}) \geq 0.$$

# What are we going to do?

• For the optimized problem (called primal in this context) we derive a dual optimization problem.

# What is it good for?

- **Simplifies optimization**. The dual problem can be of lesser complexity; in some cases the primal solution can be easily obtained from the dual solution.
- New insight. The dual problem can bring a new insight to the problem (e.g. Max-flow/Min-cut problems from graph theory are dual, or Maximum-likelihood/Minimum-entropy density estimation problems are dual).

Primal optimization problem in standard form

minimize 
$$f_0(m{x})$$
 subject to  $f_i(m{x}) \leq 0$  ,  $i=1,\ldots,m$   $h_j(m{x}) = 0$  ,  $j=1,\ldots,p$ 

where  $\mathcal{D}$  is the problem domain,  $p^*$  is the optimal value.

**Lagrangian:**  $L: \mathbb{R}^n \times \mathbb{R}^m \times \mathbb{R}^p \to \mathbb{R}$  with domain  $\operatorname{dom} L = \mathcal{D} \times \mathbb{R}^m \times \mathbb{R}^p$ 

$$L(\boldsymbol{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) = f_0(\boldsymbol{x}) + \sum_{i=1}^{m} \lambda_i f_i(\boldsymbol{x}) + \sum_{i=1}^{p} h_i \nu_i$$

- sum of objective function plus weighted sum of constraint functions
- $\lambda_i$  is Lagrange multiplier associated with  $f_i(\boldsymbol{x}) \leq 0$
- $\nu_i$  is Lagrange multiplier associated with  $h_i(\boldsymbol{x}) = 0$

# Lagrange dual function $g: \mathbb{R}^m \times \mathbb{R}^p \to \mathbb{R}$

$$g(\boldsymbol{\lambda}, \boldsymbol{\nu}) = \inf_{\boldsymbol{x} \in \mathcal{D}} L(\boldsymbol{x}, \boldsymbol{\lambda}, \boldsymbol{\nu})$$

$$= \inf_{\boldsymbol{x} \in \mathcal{D}} \left( f_0(\boldsymbol{x}) + \sum_{i=1}^m \lambda_i f_i(\boldsymbol{x}) + \sum_{i=1}^p \nu_i h_i(\boldsymbol{x}) \right)$$

- $g(\lambda, \nu)$  is a **concave function** since it is point-wise infimum of convex functions of  $(\lambda, \nu)$ ; note that it holds in general even for non-convex primal problems.
- For many important problem  $g(\lambda, \nu)$  has an analytical form.

We start form the primal LP problem

minimize 
$$oldsymbol{c}^T oldsymbol{x}$$
 subject to  $oldsymbol{A} oldsymbol{x} = oldsymbol{b}$   $oldsymbol{D} oldsymbol{x} \leq oldsymbol{q}$ 

We form the Lagrangian (using matrix notation for brevity)

$$egin{array}{lll} L(oldsymbol{x},oldsymbol{\lambda},oldsymbol{
u}) &=& f_0(oldsymbol{x}) + \sum_{i=1}^m \lambda_i f_i(oldsymbol{x}) + \sum_{i=1}^p h_i 
u_i \ &=& oldsymbol{c}^T oldsymbol{x} + oldsymbol{\lambda}^T (oldsymbol{\mathrm{D}} oldsymbol{x} - oldsymbol{q}) + oldsymbol{
u}^T (oldsymbol{\mathrm{A}} oldsymbol{x} - oldsymbol{b}) \ &=& (oldsymbol{c} + oldsymbol{\mathrm{D}}^T oldsymbol{\lambda} + oldsymbol{\mathrm{A}}^T oldsymbol{
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u}^T oldsymbol{b} \ &=& (oldsymbol{c} + oldsymbol{\mathrm{D}}^T oldsymbol{\lambda} + oldsymbol{\mathrm{A}}^T oldsymbol{
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u})^T oldsymbol{n} - oldsymbol{
u}^T oldsymbol{
u} - oldsymbo$$

We get the Lagrange dual function by minimizing w.r.t primal variables

$$g(\boldsymbol{\lambda}, \boldsymbol{\nu}) = \inf_{\boldsymbol{x}} L(\boldsymbol{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) = \left\{ \begin{array}{ll} -\boldsymbol{\lambda}^T \boldsymbol{q} - \boldsymbol{\nu}^T \boldsymbol{b} & \text{if} \quad \boldsymbol{c} + \mathbf{D}^T \boldsymbol{\lambda} + \mathbf{A}^T \boldsymbol{\nu} = \mathbf{0} \\ -\infty & \text{otherwise} \end{array} \right.$$

Weak duality: If  $\lambda \geq 0$  and  $x \in \mathcal{X}_{feas}$  then  $f_0(x) \geq g(\lambda, \nu)$ , i.e. the Lagrange dual function is a lower bound on the primal objective. In particular, it lower bounds the optimal value  $p^* \geq g(\lambda, \nu)$ ,  $\forall \lambda \geq 0$ ,  $\forall \nu$ .

To see this recall the Lagrangian

$$L(\boldsymbol{x}, \boldsymbol{\lambda}, \boldsymbol{
u}) = f_0(\boldsymbol{x}) + \sum_{i=1}^m \lambda_i f_i(\boldsymbol{x}) + \sum_{i=1}^p h_i \nu_i$$

and notice that for  $x \in \mathcal{X}_{\mathrm{feas}}$  we have:

- 1.  $f_i(\boldsymbol{x}) \leq 0$  and thus  $\sum_i \lambda_i f(\boldsymbol{x}) \leq 0$  since  $\lambda_i \geq 0$ ,
- 2.  $h_i(\boldsymbol{x}) = 0$  and thus  $\sum_i \nu_i h_i(\boldsymbol{x}) = 0$ ,

therefore

$$f_0(\boldsymbol{x}) \geq f_0(\boldsymbol{x}) + \sum_{i=1}^m \lambda_i f_i(\boldsymbol{x}) + \sum_{i=1}^p h_i \nu_i = L(\boldsymbol{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) \geq \inf_{\boldsymbol{x} \in \mathcal{D}} L(\boldsymbol{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}).$$

Note that the **weak duality holds in general** regardless the primal problem is convex or not.

# **Dual problem**

maximize  $g(\lambda, \nu)$  subject to  $\lambda \geq 0$ 

where we optimize w.r.t  $\lambda \in \mathbb{R}^m$ ,  $\nu \in \mathbb{R}^p$ ; the optimal value denoted by  $d^*$ .

- Solving the dual problem  $\approx$  finding the best lower bound  $d^*$  on primal optimal value  $p^*$  which can be obtained from the Lagrangian.
- **Duality gap** is the difference between the primal and the dual optimal values  $p^* d^* \ge 0$ , i.e., it determines the tightness of the lower bound.
- The dual problem is always convex since  $g(\lambda, \nu)$  is a concave function regardless the primal problem is convex or not.
- $(\lambda, \nu)$  are **dual feasible** if  $\lambda \geq 0$  and  $g(\lambda, \nu) > -\inf$ , i.e. for dual feasible points we have non-trivial lower bound.

It usually helps if the constraint  $g(\lambda, \nu) > -\inf$  is expressed explicitly in the dual problem.

# The primal LP problem

minimize 
$$oldsymbol{c}^T oldsymbol{x}$$
 subject to  $oldsymbol{A} oldsymbol{x} = oldsymbol{b}$   $oldsymbol{D} oldsymbol{x} \leq oldsymbol{q}$ 

with the Lagrange dual function

$$g(\boldsymbol{\lambda}, \boldsymbol{\nu}) = \begin{cases} -\boldsymbol{\lambda}^T \boldsymbol{q} - \boldsymbol{\nu}^T \boldsymbol{b} & \text{if} \quad \boldsymbol{c} + \mathbf{D}^T \boldsymbol{\lambda} + \mathbf{A}^T \boldsymbol{\nu} = \mathbf{0} \\ -\infty & \text{otherwise} \end{cases}$$

The dual problem reads

maximize 
$$g(\lambda, \nu)$$
 subject to  $\lambda \geq 0$ 

Making the constraint  $g(\lambda, \nu) > -\inf$  explicit, i.e.,  $c + \mathbf{D}^T \lambda + \mathbf{A}^T \nu = \mathbf{0}$ , we get the **dual LP problem** 

maximize 
$$- \pmb{\lambda}^T \pmb{q} - \pmb{\nu}^T \pmb{b}$$
 subject to  $\pmb{\lambda} \geq 0$   $\pmb{c} + \mathbf{D}^T \pmb{\lambda} + \mathbf{A}^T \pmb{\nu} = \pmb{0}$ 

**Strong duality** holds if the duality gap is zero, i.e.,  $p^* = d^*$  and the Lagrangian lower bound is tight.

# When does it happen?

- It does not hold in general.
- It holds if the primal problem is convex and the Slater's condition (also called constraint qualification) holds:
  - **Slater's condition** holds if there exists a strictly feasible point, i.e., there exists  $x \in \mathcal{X}_{\text{feas}}$  such that  $f_i(x) < 0$ ,  $i = 1, \ldots, m$ ; note that this condition is very mild.
- There also exist non-convex problems for which the strong duality holds.

A triplet  $(x, \lambda, \nu)$  satisfy the Karush-Kuhn-Tucker conditions if:

$$rac{\partial L(oldsymbol{x},oldsymbol{\lambda},oldsymbol{
u})}{\partial oldsymbol{x}} = oldsymbol{0}$$

partial derivative of L w.r.t  $\boldsymbol{x}$  vanishes

$$\frac{\partial L(\boldsymbol{x}, \boldsymbol{\lambda}, \boldsymbol{\nu})}{\partial \boldsymbol{\lambda}} \leq \mathbf{0}$$

implies  $f_i(\boldsymbol{x}) \leq 0$ ,  $i = 1, \dots, m$ 

$$\frac{\partial L(\boldsymbol{x}, \boldsymbol{\lambda}, \boldsymbol{\nu})}{\partial \boldsymbol{\nu}} = \mathbf{0}$$

implies  $h_i(\boldsymbol{x}) = 0$ ,  $i = 1, \dots, p$ .

$$\lambda \geq 0$$

duality constraint holds

$$\lambda_i f_i(\boldsymbol{x}) = 0$$
,  $i = 1, \dots, m$ 

 $\lambda_i f_i(\boldsymbol{x}) = 0$ ,  $i = 1, \dots, m$  so called complementary slackness

- If strong duality holds then KKT conditions are necessary for  $(\boldsymbol{x}, \boldsymbol{\lambda}, \boldsymbol{\nu})$  to be optimal.
- If primal problem is convex and Slater's condition holds then KKT conditions are necessary and sufficient for  $(x, \lambda, \nu)$  to be optimal.

The primal LP problem

minimize 
$$oldsymbol{c}^T oldsymbol{x}$$
 subject to  $oldsymbol{A} oldsymbol{x} = oldsymbol{b}$   $oldsymbol{D} oldsymbol{x} \leq oldsymbol{q}$ 

with the Lagrangian

$$L(\boldsymbol{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) = \boldsymbol{c}^T \boldsymbol{x} + \boldsymbol{\lambda}^T (\mathbf{D} \boldsymbol{x} - \boldsymbol{q}) + \boldsymbol{\nu}^T (\mathbf{A} \boldsymbol{x} - \boldsymbol{b})$$

The KKT conditions read:

$$egin{array}{lll} rac{\partial L(oldsymbol{x},oldsymbol{\lambda},oldsymbol{
u})}{\partial oldsymbol{x}} = oldsymbol{0} &\Rightarrow oldsymbol{c} + \mathbf{D}^Toldsymbol{\lambda} + \mathbf{A}^Toldsymbol{
u} = oldsymbol{0} \ rac{\partial L(oldsymbol{x},oldsymbol{\lambda},oldsymbol{
u})}{\partial oldsymbol{\lambda}} \leq oldsymbol{0} &\Rightarrow oldsymbol{D}oldsymbol{x} - oldsymbol{q} \leq oldsymbol{0} \ rac{\partial L(oldsymbol{x},oldsymbol{\lambda},oldsymbol{
u})}{\partial oldsymbol{
u}} = oldsymbol{0} &\Rightarrow oldsymbol{D}oldsymbol{x} - oldsymbol{q} = oldsymbol{0} \ rac{\partial L(oldsymbol{x},oldsymbol{
u},oldsymbol{
u})}{\partial oldsymbol{
u}} = oldsymbol{0} &\Rightarrow oldsymbol{A}oldsymbol{x} - oldsymbol{b} = oldsymbol{0} \ \lambda \geq oldsymbol{0} \ \lambda_i f_i(oldsymbol{x}) = oldsymbol{0}, \ i = 1, \ldots, m &\Rightarrow oldsymbol{\lambda}^T(oldsymbol{D}oldsymbol{x} - oldsymbol{q}) = oldsymbol{0} \ \end{pmatrix}$$

Let us consider an unconstrained convex problem

minimize 
$$f(x)$$

### **General descent method:**

Initialization: set  $x \in \text{dom } f$ .

# repeat

- 1. Determine a descent direction  $\delta$ .
- 2. Line-search: find a step size  $t = \operatorname{argmin}_{t'>0} f(\boldsymbol{x} + t'\boldsymbol{\delta})$ .
- 3. Update  $x := x + t\delta$ .

until stopping condition is satisfied.

- ullet It generates a sequence of  $oldsymbol{x}^{(1)}, oldsymbol{x}^{(2)}, \dots$  such that  $f(oldsymbol{x}^{(k)}) > f(oldsymbol{x}^{(k+1)}).$
- ullet For f differentiable, a vector  $oldsymbol{\delta}$  is a descent direction if

$$f'(\boldsymbol{x}; \boldsymbol{\delta}) = \lim_{h \to 0_+} \frac{f(\boldsymbol{x} + h\boldsymbol{\delta})}{h} = \nabla f(\boldsymbol{x})^T \boldsymbol{\delta} < 0$$

e.g., gradient descent methods use  $\delta = -\nabla f(x)$ .

Let us consider equality constrained convex problem

$$\begin{array}{ll} \text{minimize} & f(\boldsymbol{x}) \\ \text{subject to} & \mathbf{A}\boldsymbol{x} = \boldsymbol{b} \end{array}$$

ullet Using the KKT optimality conditions,  $m{x} \in \mathbf{dom}\, f$  is optimal iff there exist  $m{
u}$  such that

$$\mathbf{A}\mathbf{x} = \mathbf{b}$$
,  $\nabla f(\mathbf{x}) + \mathbf{A}^T \mathbf{\nu} = 0$ .

• For a convex quadratic function  $f(x) = \frac{1}{2}x^T\mathbf{H}x + c^Tx$  the KKT conditions lead to an efficiently solvable set of linear equations:

$$\mathbf{A}x = \mathbf{b}$$
,  $\mathbf{H}x + \mathbf{c} + \mathbf{A}^T \mathbf{\nu} = 0$ .

• Newton method is applicable for a general twice differentiable function f(x): it iteratively approximates f(x) by a quadratic function

$$\hat{f}(x) = \frac{1}{2}(x - x')\nabla^2 f(x')(x - x') + \nabla f(x')^T(x - x') + f(x')$$

and solves the KKT conditions for the approximation  $\hat{f}(x)$ .

Let us consider equality constrained convex problem

minimize 
$$f_0(m{x})$$
 subject to  $f_i(m{x}) \leq 0$  ,  $i=1,\ldots,m$   $\mathbf{A}m{x} = m{b}$ 

ullet Constraints  $f_i(oldsymbol{x}) \leq 0$  can be made implicit using the **barrier function** 

$$\phi_i(\boldsymbol{x}) = \begin{cases} 0 & \text{if } f_i(\boldsymbol{x}) \leq 0 \\ \infty & \text{if } f_i(\boldsymbol{x}) > 0 \end{cases}$$

i.e., we can equivalently optimized equality constraint problem

minimize 
$$f_0(\boldsymbol{x}) + \sum_{i=1}^m \phi_i(\boldsymbol{x})$$
 subject to  $\mathbf{A}\boldsymbol{x} = \boldsymbol{b}$ 

ullet Functions  $\phi_i(oldsymbol{x})$  are approximated by a differentiable convex functions

$$\hat{\phi}_i(\boldsymbol{x}) = -\frac{1}{t}\log(-f_i(\boldsymbol{x})),$$

which for high t well approximates the step barrier function  $\phi_i(x)$ .

# Materials used to prepare this lecture:

• S. Boyd, L. Vandenberghe: *Convex optimization*. Cambridge University Press. 2004.

Available at: http://www.stanford.edu/~boyd/cvxbook/

- S. Boyd: Lecture notes for EE364, Stanford University. 2007-2008. Available at: http://www.stanford.edu/class/ee364/
- H. Hindi: A Tutorial on Convex Optimization II: Duality and Interior Point Methods. Palo Alto Research Center, California.

Google: hindi tutorial convex

### **Further recommended literature:**

- D.P. Bertsekas. *Nonlinear Programming*. (2nd edition), Athena Scientific, Belmont, Massachusetts, 1999.
- J.F. Bonnans, et. al: *Numerical Optimization*. (2nd edition), Springer, Heidelberg, 2006.